## MATH 245 F23, Exam 1 Solutions

1. Carefully define the following terms: composite, converse.

Let $n$ be an integer with $n \geq 2$. We call $n$ composite if there exists some integer $m$ satisfying both $m \mid n$ and $1<m<n$. For arbitrary propositions $p, q$, the converse of conditional proposition $p \rightarrow q$ is the proposition $q \rightarrow p$.
2. Carefully state the following theorems: Division Algorithm Theorem, Disjunctive Syllogism Theorem.
The Division Algorithm Theorem says: For any integers $a, b$ with $b \geq 1$, there are unique integers $q, r$ satisfying $a=b q+r$ and $0 \leq r<b$. The Disjunctive Syllogism Theorem states: for any propositions $p, q$, if $p \vee q$ is $T$ and $q$ is F , then $p$ must be $T$.
3. Use a truth table to help prove the following:

For all propositions $p, q$, we have $(p \uparrow q),(p \rightarrow q) \vdash \neg p$.
Let $p, q$ be arbitrary propositions. Suppose that $(p \uparrow q),(p \rightarrow q)$ are both $T$. Consider the truth table at right. Because $p \uparrow q$ is $T$, the first row is impossible. Since $p \rightarrow q$ is $T$, the second row is impossible. In both

| $p$ | $q$ | $p \uparrow q$ | $p \rightarrow q$ | $\neg p$ |
| :---: | :---: | :---: | :---: | :---: |
| T | T | F | T | F |
| T | F | T | F | F |
| F | T | T | T | T |
| F | F | T | T | T | remaining rows, $\neg p$ is $T$.

4. Let $p, q$ be propositions. Without using truth tables, prove $p \wedge q \equiv q \wedge p$.

Note: do not use/cite commutativity of $\wedge-$ you are being asked to prove commutativity! There are four cases, but three of them end up collapsing. If $p, q$ are both $T$, then $p \wedge q$ is $T$, but also $q \wedge p$ is $T$ since $q, p$ are both $T$. If $p, q$ are not both $T$ (the three cases of $p, q$ both F , or $p$ is F and $q$ is T , or $p$ is T and $q$ is F ), then $p \wedge q$ is F , but also $q \wedge p$ is $F$ since $q, p$ are not both $T$. In all cases, $p \wedge q$ agrees with $q \wedge p$.
ALTERNATE SOLUTION: Suppose $p \wedge q$ is $T$. By simplification twice, $p$ and $q$ are $T$. By conjunction, $q \wedge p$ is $T$. This proves $p \wedge q \vdash q \wedge p$. Now suppose $q \wedge p$ is $T$. By simplification twice, $q$ and $p$ are $T$. By conjunction, $p \wedge q$ is $T$. This proves $q \wedge p \vdash p \wedge q$. Together these prove $p \wedge q \equiv q \wedge p$.
5. Let $p, q$ be propositions. Use semantic theorems to prove the "Trivial Proof Theorem": $q \vdash p \rightarrow q$. Do not use the theorem to prove itself!
We begin by assuming $q$. By addition, $q \vee(\neg p)$. By conditional interpretation, $p \rightarrow q$.
6. Prove or disprove: For all $a, b, c \in \mathbb{Z}$, if $a c \mid b$ then $a \mid b$.

The statement is true. We begin by letting $a, b, c \in \mathbb{Z}$ be arbitrary, and apply a direct proof. Suppose that $a c \mid b$. Then there is some integer $k$ with $a c k=b$. Set $m=c k$, which is an integer since $c, k$ are. We have $a m=b$ for an integer $m$, so $a \mid b$.
7. Prove or disprove: For all $x \in \mathbb{Z},|4 x+9|>1$.

The statement is false, so we need a counterexample. Take $x^{\star}=-2$. We have $\left|4 x^{\star}+9\right|=|4(-2)+9|=|1|=1 \ngtr 1$.
8. Let $a, b \in \mathbb{N}_{0}$. Use the definition of $\leq$ to prove that if $2 a \leq b$ then $2 a^{2}+a b \leq b^{2}$. We will use a direct proof. Suppose that $2 a \leq b$. Then $b-2 a \in \mathbb{N}_{0}$. Also $b+a \in \mathbb{N}_{0}$, since $a, b \in \mathbb{N}_{0}$. But now the product $(b-2 a)(b+a) \in \mathbb{N}_{0}$, i.e. $b^{2}-a b-2 a^{2} \in \mathbb{N}_{0}$. Hence $b^{2}-\left(2 a^{2}+a b\right) \in \mathbb{N}_{0}$, so $2 a^{2}+a b \leq b^{2}$.

Note: How would someone come up with the strange idea that $a+b \in \mathbb{N}_{0}$, and we can multiply by it? Work backwards from the end, and factor $b^{2}-2 a^{2}-a b$ (which is easy to do if we know that $b-2 a$ is a factor).
9. Prove or disprove: For all $a, b, c \in \mathbb{Z}$, if $a \mid b$, then $a c^{2} \mid b^{2} c$.

The statement is false, so we need a counterexample. To disprove an implication $p \rightarrow q$ we need to make $p$ true and $q$ false, i.e. we need our example to satisfy $a \mid b$ and $a c^{2} \nmid b^{2} c$. Many counterexamples are possible.

One choice is $a=2, b=10, c=3$. We have $2 \mid 10$ since $2 \cdot 5=10$. This proves that $a \mid b$. Now, $a c^{2}=18$ and $b^{2} c=300$. If $18 k=300$, then $k=\frac{300}{18}=\frac{50}{3}=16 \frac{2}{3}$, which is not an integer. Hence $a c^{2} \nmid b^{2} c$.
Perhaps the simplest choice is $a=b=1, c=2$, although some students don't like simple choices like this. We have $1 \mid 1$ since $1 \cdot 1=1$. This proves that $a \mid b$. Now, $a c^{2}=4$ and $b^{2} c=2$. If $4 k=2$, then $k=0.5 \notin \mathbb{Z}$. Hence $a c^{2} \nmid b^{c}$.
10. Prove or disprove: $\forall x, y \in \mathbb{R},(x<y) \rightarrow(\exists z \in \mathbb{R}, x<z<y)$.

The statement is true. We begin by letting $x, y \in \mathbb{R}$ be arbitrary. Via direct proof, we assume that $x<y$.
We now need to prove $\exists z \in \mathbb{R}, x<z<y$; it's a little tricky to find such a $z$. The usual method is to take the midpoint, i.e. take $z=\frac{x+y}{2}=\frac{x}{2}+\frac{y}{2}$. Now since $x<y$ we have $\frac{x}{2}<\frac{y}{2}$. Adding $\frac{x}{2}$ to both sides gives $x<z$, while adding $\frac{y}{2}$ to both sides gives $z<y$. Combining these, we get $x<z<y$.

